

Techniques de réécriture

TD n°1 : Proofs of termination *

Exercise 1 :

Are those rewrite systems terminating?

- $\{f(a) \rightarrow f(b), g(b) \rightarrow g(a)\}$
- $\{f(f(x)) \rightarrow f(g(f(x)))\}$
- $\{f(g(x)) \rightarrow g(f(x))\}$
- $\{f(g(g(f(x)))) \rightarrow g(f(f(g(x))))\}$
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- $\{f(g(g(x))) \rightarrow g(f(f(g(x))))\}$

Exercise 2 :

- 1) Prove that $R_1 = \{or(x, y) \rightarrow x, or(x, y) \rightarrow y\}$ terminates.
- 2) Prove that $R_2 = \{f(a, b, x) \rightarrow f(x, x, x)\}$ terminates.
- 3) What can you say about $R_1 \cup R_2$?

Exercise 3 :

Prove that those functions terminate using rewrite systems :

let rec Ack = function

(0, x) -> x + 1
| (x, 0) -> Ack(x - 1, 1)
| (x, y) -> Ack(x - 1, Ack(x, y - 1));;

let rec f = function

(0, x) -> 1
| (x, 0) -> 1
| (x, y) -> f(x - 1, x - 1) + 2 * f(x - 1, y - 1) + f(y - 1, y - 1);;

Exercise 4 :

Heracles vs Lernaean Hydra. The Hydra is a term on the alphabet $\mathcal{F} = \{n(*), h(0)\}$. When Heracles cuts a head, an arbitrary number of heads grows back this way :

$$n(x_1, \dots, x_p, n(y_1, \dots, y_q, h), z_1, \dots, z_r) \rightarrow n(x_1, \dots, x_p, n(y_1, \dots, y_q), h, \dots, h, z_1, \dots, z_r)$$

Heracles wins when there is no head to cut i.e. the term representing the Hydra is irreducible.

Prove that Heracles always wins i.e. this rewrite system is terminating.

Exercise 5 :

We say that a rewrite system \mathcal{R} :

- goes around in circles when $\exists t, t \rightarrow_{\mathcal{R}}^+ t$
- loops when $\exists t, t', t \rightarrow_{\mathcal{R}}^+ t'$ and t is a sub-term of t'
- is self-imbricated when $\exists t, t', t \rightarrow_{\mathcal{R}}^+ t'$ and $t \leq t'$
- weakly terminates when $\forall t, \exists t', t \rightarrow_{\mathcal{R}}^+ t'$ and t' irreducible

- 1) Prove the following implications :

goes around in circles \Rightarrow loops \Rightarrow does not terminate \Rightarrow is self-imbricated

- 2) Give a counter-example for each inverse implication.

- 3) Give an example of a rewrite system which weakly terminates but does not terminate.

Exercise 6 :

For which values of n, k, m and p does the rewrite system $\{f^n(g^k(x)) \rightarrow g^m(f^p(x))\}$ terminate?

*taken from *Les termes en logique et en programmation*