

Techniques de réécriture

TD n°3 : RPO & KBO

Exercise 1 :

Which rewrite systems from TD 1 are terminating using RPO ?

Definition. We recall that a strict order $>$ on $T(\mathcal{F}, \mathcal{X})$ is a simplification order if :
compatibility : $\forall s_1 > s_2, t_1, \dots, t_{i-1}, t_{i+1}, \dots, t_n$ and $f \in \mathcal{F}$,

$$f(t_1, \dots, t_{i-1}, s_1, t_{i+1}, \dots, t_n) > f(t_1, \dots, t_{i-1}, s_2, t_{i+1}, \dots, t_n)$$

substitution closure : $\forall s_1 > s_2, \sigma$ substitution, $s_1\sigma > s_2\sigma$

subterm property : $\forall t, p \in Pos(t) \setminus \{\epsilon\}, t > t|_p$

Definition. Fix $>$ a strict order on \mathcal{F} . A weight function is a function $w : \mathcal{F} \cup \mathcal{X} \rightarrow \mathbb{R}_+$.

We say that w is admissible w.r.t. $>$ if :

— $\exists w_0 \in \mathbb{R}_+^*$ such that $\forall x \in \mathcal{X}, w(x) = w_0$ and for all constant $c, w(c) \geq w_0$

— if $f \in \mathcal{F}$ is of arity 1 and $w(f) = 0$ then f is maximal for $>$

We extend a weight function on $T(\mathcal{F}, \mathcal{X})$ by : $w(t) = \sum_{x \in \mathcal{X}} w(x). |t|_x + \sum_{f \in \mathcal{F}} w(f). |t|_f$.

Definition. Fix a strict order $>$ on \mathcal{F} and a weight function w admissible w.r.t. $>$. The Knuth-Bendix order $<_{kbo}$ induced on $T(\mathcal{F}, \mathcal{X})$ is defined by $s <_{kbo} t$ if $\forall x \in \mathcal{X}, |s|_x \leq |t|_x$ and if one of the following holds :

(KBO1) $w(s) < w(t)$

(KBO2) $w(s) = w(t)$ and one of the following holds

(KBO2a) $\exists f \in \mathcal{F}, x \in \mathcal{X}$ and $k > 0$ such that $s = x$ and $t = f^k(x)$

(KBO2b) $\exists f > g$ such that $s = g(s_1, \dots, s_m)$ and $t = f(t_1, \dots, t_n)$

(KBO2c) $\exists f \in \mathcal{F}$ and $i \leq n$ such that $s = f(s_1, \dots, s_n), t = f(t_1, \dots, t_n), s_1 = t_1, \dots, s_{i-1} = t_{i-1}$ and $s_i <_{kbo} t_i$

Exercise 2 :

Let $R = \{h(f(x), y) \rightarrow f(g(x), y), g(x), y) \rightarrow h(x, y)\}$.

1) Show that you cannot prove that R is terminating using RPO.

2) Prove it is terminating using an extension of a RPO.

3) Assuming that $<_{kbo}$ is a simplification order, prove that R is terminating.

Exercise 3 :

1) Prove that $<_{kbo}$ is compatible, closed under substitutions and has the subterm property.

2) Prove that $<_{kbo}$ is a strict ordering (irreflexive and transitive). Conclude.

Exercise 4 :

Which rewrite systems from TD 1 are terminating using KBO ?